

① Euclidean Algorithm:

$$\begin{array}{ll} \text{line 1} & a = bq_1 + r_1 \quad 0 < r_1 < b \\ \text{line 2} & b = r_1q_2 + r_2 \quad 0 < r_2 < r_1 \\ \text{line 3} & r_1 = r_2q_3 + r_3 \quad 0 < r_3 < r_2 \\ & \vdots \\ \text{line } k-1 & r_{k-3} = r_{k-2}q_{k-1} + r_{k-1} \quad 0 < r_{k-1} < r_{k-2} \\ \text{line } k & r_{k-2} = r_{k-1}q_k + r_k \quad 0 < r_k < r_{k-1} \\ \text{line } k+1 & r_{k-1} = r_kq_{k+1} + 0 \end{array}$$

To Prove) $\leftarrow$ N.T.P two things to prove $r_k = \gcd(a, b)$

① $r_k \mid a$ and $r_k \mid b$

② if $r \mid a$ and $r \mid b$ then $r \leq r_k$

Proof ① Line $k+1 \Rightarrow r_k \mid r_{k-1}$

Then since $r_{k-2} = r_{k-1}q_k + r_k$ (line k)
 r_k also divides r_{k-2} .

Then line $k-1 \Rightarrow r_k \mid r_{k-3}$

Continuing the process we get $r_k \mid b$

then $r_k \mid a$.

② Suppose $r \mid a$ and $r \mid b$.

Then $r \mid r_1$ because line 1 $\Rightarrow r_1 = a - bq_1$
 then similarly b/c $r \mid b$ and $r \mid r_1$ we get $r \mid r_2$
 (because line 2 $\Rightarrow r_2 = b - r_1q_2$)

Continuing this process we get that $r \mid r_k$ so $r \leq r_k$

By ① & ②, $r_k = \gcd(a, b)$.

② Euclid's lemma

② p prime, $p \mid ab \Rightarrow p \mid a$ or $p \mid b$

Proof Suppose p prime $p \mid ab$ and $p \nmid a$

need to show $\leftarrow$ N.T.S. $p \mid b$.

p prime, $p \nmid a \Rightarrow \gcd(a, p) = 1$

$\Rightarrow \exists s, t \in \mathbb{Z}$ such that $1 = as + pt$

Then $b = bas + bpt$

Since $p \mid ab$ we have $p \mid bas$.

Obviously $p \mid bpt$

So, $p \mid bas + bpt \Rightarrow p \mid b$.

③ Ch. 0 #4

let $a = 7$, $b = 11$

$$1 = 3 \cdot 7 - 2 \cdot 11$$

also

$$1 = 8 \cdot 7 - 5 \cdot 11$$

Ch 0 #8

$a, b, c \in \mathbb{Z}$ s.t. $a \mid c$ and $b \mid c$

If $\gcd(a, b) = 1$ show that $ab \mid c$

Show if $\gcd(a, b) \neq 1$ then result doesn't hold.

Proof: Suppose $a \mid c$, $b \mid c$, $\gcd(a, b) = 1$ and $ab \nmid c$.

Then $\exists$ a prime p s.t. $p \mid ab$ but $p \nmid c$. Since $\gcd(a, b) = 1$ by Euclid's lemma either $p \mid a$ or $p \mid b$.

If $p \mid a$, since $p \nmid c$ we get $a \nmid c$.

If $p \mid b$ since $p \nmid c$ we get $b \nmid c$.

So, we can't have $ab \nmid c$.

Hence $ab \mid c$.

Ch 0 #16

$$126 = 3 \times 34 + 24$$

$$34 = 1 \times 24 + 10$$

$$24 = 2 \times 10 + 4$$

$$10 = 2 \times 4 + 2 \quad \text{gcd}$$

$$4 = 2 \times 2 + 0$$

$$2 = 10 - 2 \times 4$$

$$= 10 - 2 \times (24 - 2 \times 10) = 5 \times 10 - 2 \times 24$$

$$= 5 \times (34 - 24) - 2 \times 24 = 5 \times 34 - 7 \times 24$$

$$= 5 \times 34 - 7 \times (126 - 3 \times 34)$$

$$2 = 26 \times 34 - 7 \times 126$$

Ch 0 #21

A set with n elements has 2^n subsets
for every $n \in \mathbb{Z}^+$

Proof by Induction:

Initial case: $n=1$

Subsets of a set, with a single element
are $\emptyset$ and the set itself so

For $n=1$, #subsets $= 2^1 = 2$

inductive hypothesis:

Assume a set with n elements
has 2^n subsets.

Suppose A has $n+1$ elements,

Say $A = \{a_1, a_2, \dots, a_n, a_{n+1}\}$

Consider $B = \{a_1, \dots, a_n\}$

By the inductive assumption B
has 2^n subsets.

Subsets of B are obviously
subsets of A . A will have
additional subsets which are
obtained by including a_{n+1} to
each of the subsets of B .

So number of additional subsets
will be 2^n as well.

Total number of subsets of A
is then $2^n + 2^n = 2 \cdot 2^n = 2^{n+1}$.

(We used the first principle of induction) $\square$

Ch 0 #41

ISBN for our book: $\langle 0 \ 6 \ 1 \ 8 \ 5 \ 1 \ 4 \ 7 \ 16 \rangle$

dot product with $\langle 10, 9, 8, 7, 6, 5, 4, 3, 2, 1 \rangle$ is $54 + 8 + 56 + 30 + 5 + 16 + 21 + 2 + 6$

mod 11 $-1 + 8 + 2 + 7 + 5 + 5 + (-1) + 2 + 6 = 0$

(5) $S = \mathbb{Z}$, $a R b$ iff $5 | a - b$.

i) reflexivity: Is $a R a \forall a \in S$?

i.e. Does $5 | a - a \forall a \in \mathbb{Z}$? yes

ii) symmetry If $a R b$ then is it true $b R a \forall a, b \in S$?

$a R b \Rightarrow 5 | a - b \Rightarrow 5 | -(a - b) \Rightarrow 5 | b - a \Rightarrow b R a$ ✓

iii) transitivity If $a R b$ and $b R c$ then is $a R c \forall a, b, c \in S$?

$a R b$ and $b R c \Rightarrow 5 | a - b$ and $5 | b - c \Rightarrow 5 | a - b + b - c \Rightarrow 5 | a - c \Rightarrow a R c$ ✓

equivalence classes:

$$\bar{0} = \{a \in \mathbb{Z} \mid 5 | a - 0\} = \{5k \mid k \in \mathbb{Z}\}$$

$$\bar{1} = \{a \in \mathbb{Z} \mid 5 | a - 1\} = \{5k + 1 \mid k \in \mathbb{Z}\}$$

$$\bar{2} = \{5k + 2 \mid k \in \mathbb{Z}\}$$

similar for $\bar{3}$ and $\bar{4}$.

Ch 0 #28

f_n : n th Fibonacci
number

$$f_1 = f_2 = 1$$

$$f_n = f_{n-1} + f_{n-2} \text{ for } n \geq 3$$

Show that $f_n < 2^n$ th

Proof by Induction

initial case $n=3$

$$f_3 = f_2 + f_1$$

$$= 1 + 1 < 8 = 2^3$$

inductive hypothesis:

Assume $f_k < 2^k$ for all $k \leq n$.

Prove that $f_n < 2^n$. second principle

$$f_n = f_{n-1} + f_{n-2}$$

By inductive assumption

$$f_{n-1} < 2^{n-1} \text{ and } f_{n-2} < 2^{n-2}$$

$$\text{So, } f_n < 2^{n-1} + 2^{n-2}$$

$$< 2^{n-1} + 2^{n-1}$$

$$= 2 \cdot 2^{n-1} = 2^n$$

So, $f_n < 2^n$